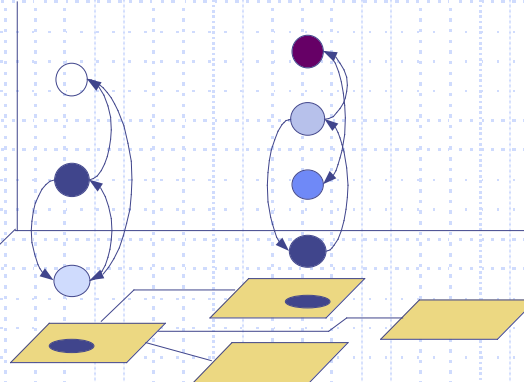


## State Machines and Object Life Cycles

G. Falquet, L. Nerima

### Another Modelling Dimension

- ❖ Dynamic view of the objects
- ❖ How objects evolve, change their state



## State Machine

- ❖ Specify a sequence of states that an object goes through in response to events
- ❖ Describe object life cycles
- ❖ Centred on one object
- ❖ Contains object states and state transitions
- ❖ Strong theoretical background: finite state automata, Petri nets, state charts (Harel). Real-time systems.

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## State (of an object)

- ❖ Situation during the life of an object
- ❖ During this situation the object
  - ❖ satisfies some condition
  - ❖ or performs some **activity**
  - ❖ or waits for some event
- ❖ A state can have substates (finer description)

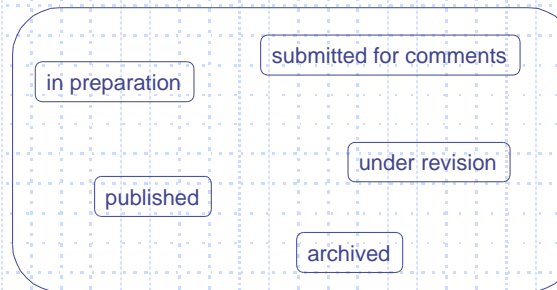
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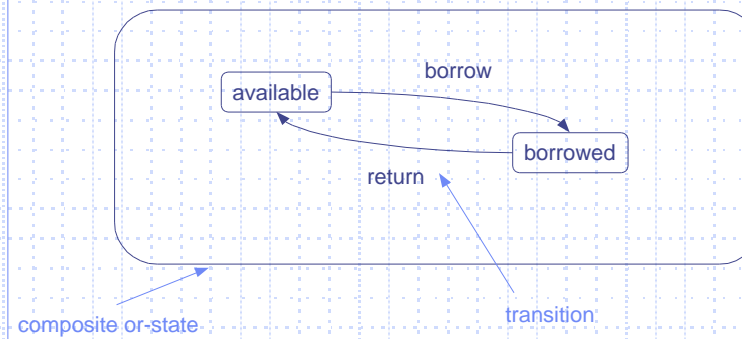
## Example: Document



## Event

- ❖ "Noteworthy occurrence that has a location in time and space."
- ❖ Can trigger a state transition (+ actions).
- ❖ Kinds of event:
  - ❖ call (operation),
  - ❖ change (something has changed in the system),
  - ❖ signal (send an event to another object),
  - ❖ time (timer, timeout, ...)

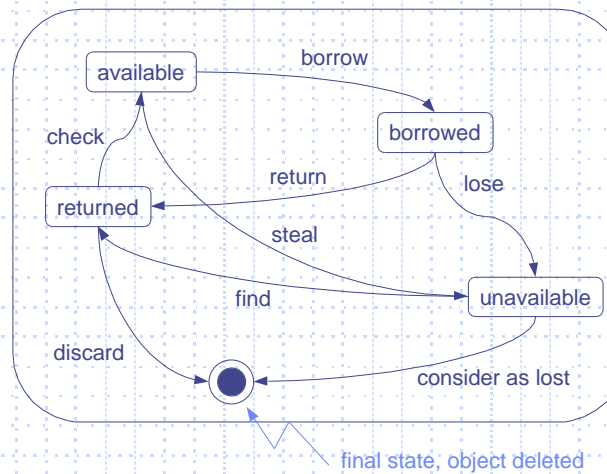
## A Simple State Machine for Books



- ❖ A transition can be fired only if the object is in the source state and the event occurs.

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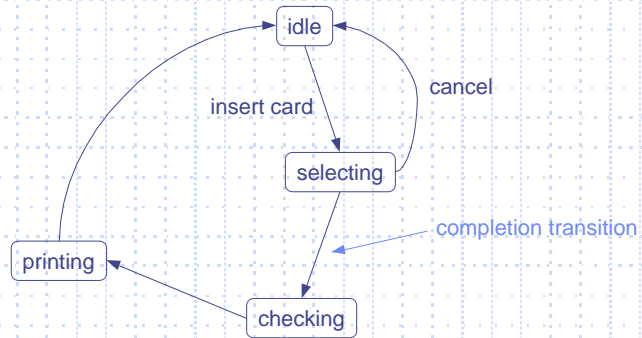
## Another State Machine for Books



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## Completion transitions

- ❖ Triggered by the completion of the state's activity
- ❖ Notation: no event name



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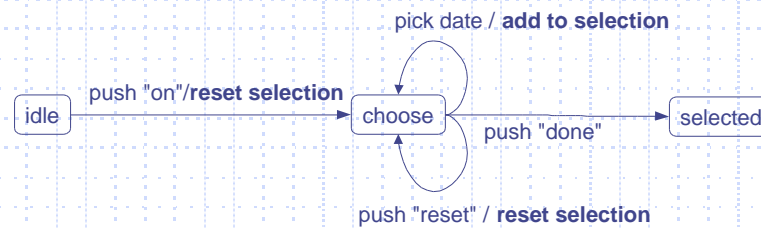
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## Event triggered actions

- ❖ A transition may execute an action
- ❖ Actions are atomic (assignment, send signal, create object, create link, ...)
- ❖ Actions must be fast ( complex computation)



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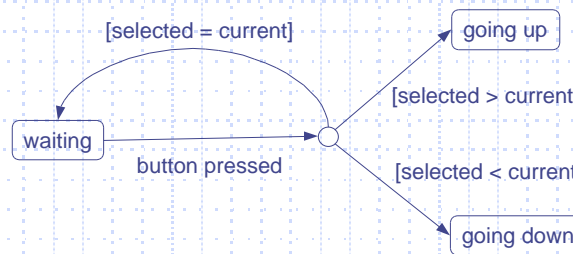
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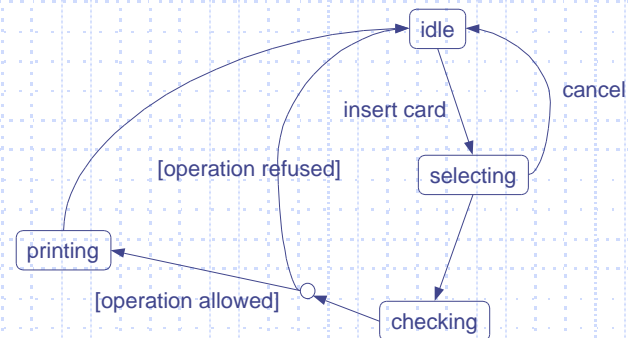
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## Guarded transitions

- ❖ Guard = additional condition on a transition
- ❖ Transition fired only if guard is true
- ❖ Only one transition fired if several guards are true
- ❖ Guards are queries (may not change values)



## Guarded completion

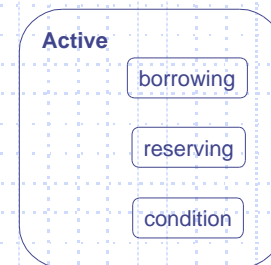


## Concurrent states

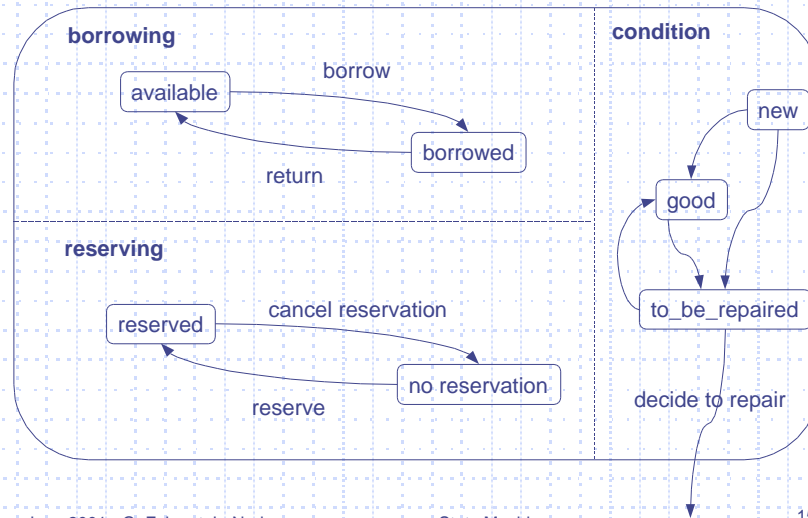
- ❖ An object may be in two (or more) states at the same time.
  - ❖ Book: (borrowed + reserved)
  - ❖ Book: (borrowed + not reserved)
- ❖ => Large number of combinations
- ❖ Composite concurrent states

## Concurrent states (*and* states)

- ❖ Being in state Active means being simultaneously in states condition, borrowing, reserving.



## Composite concurrent states

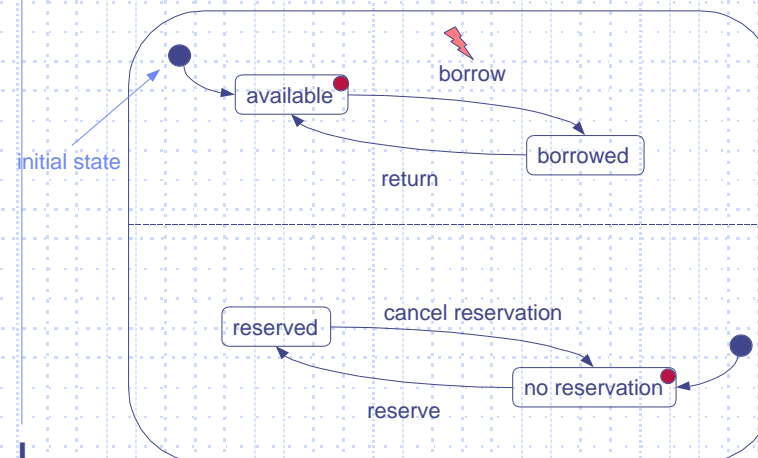


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## Concurrent active states



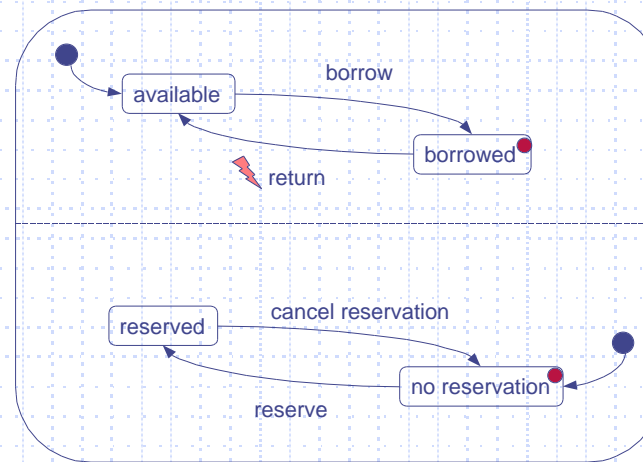
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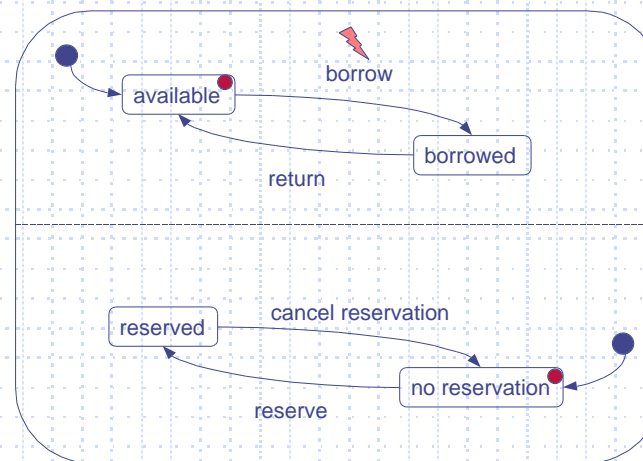
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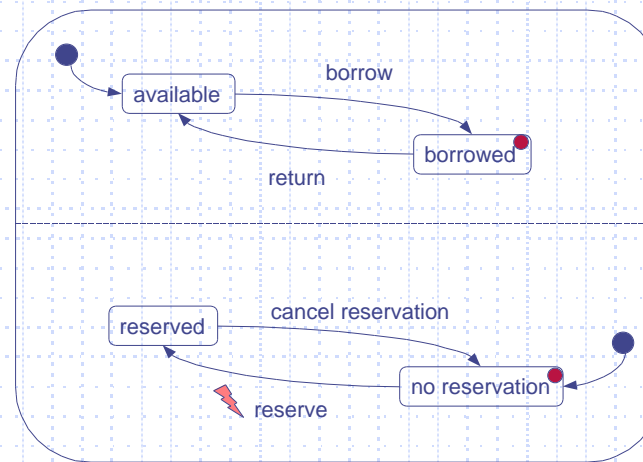
## Concurrent active states



## Concurrent active states



## Concurrent active states



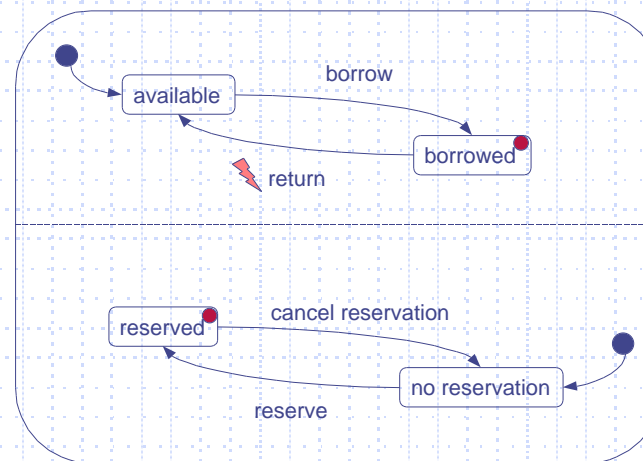
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## Concurrent active states



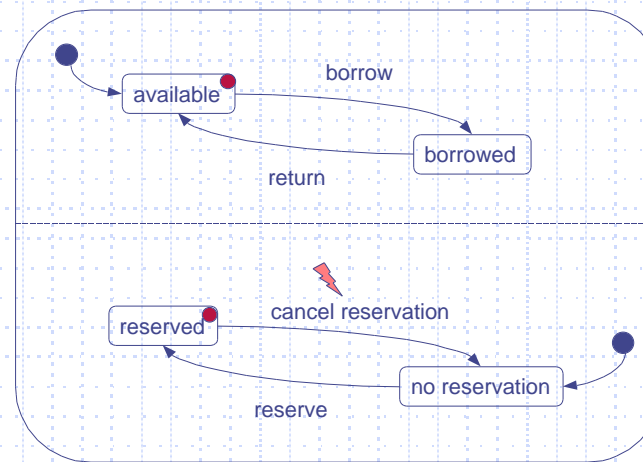
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## Concurrent active states



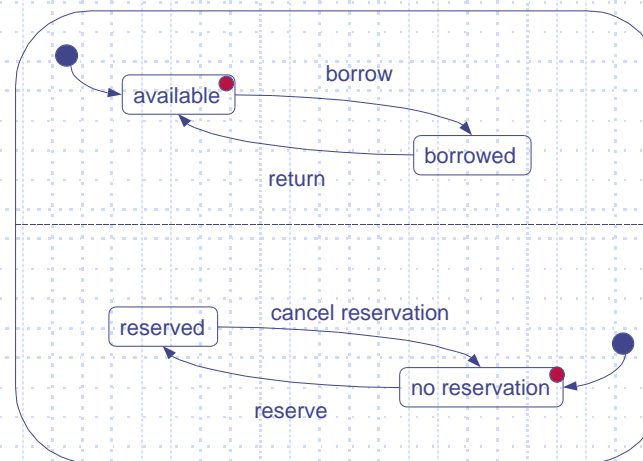
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## Concurrent active states



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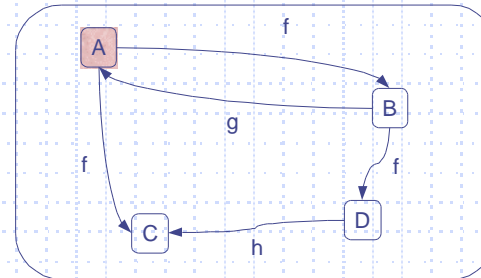
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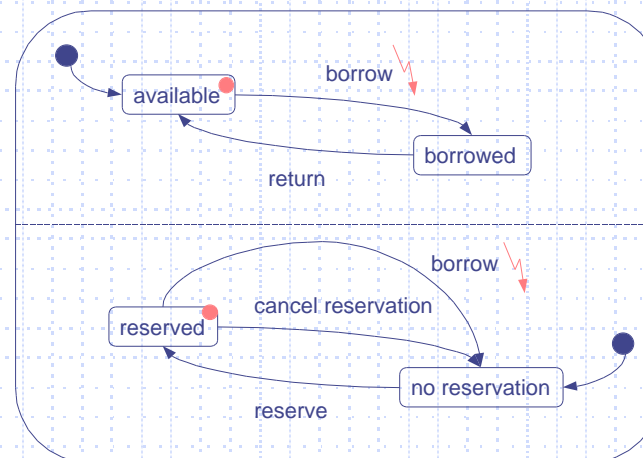
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## Transitions on the same event

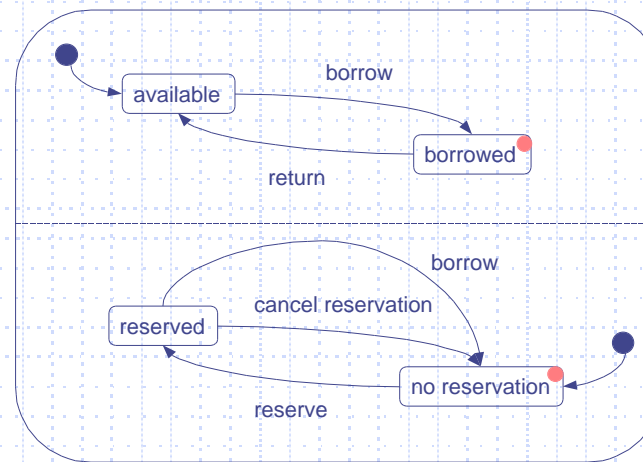
- ❖ What happens if A is active and f occurs ?
- ❖ Non determinism in an or-state



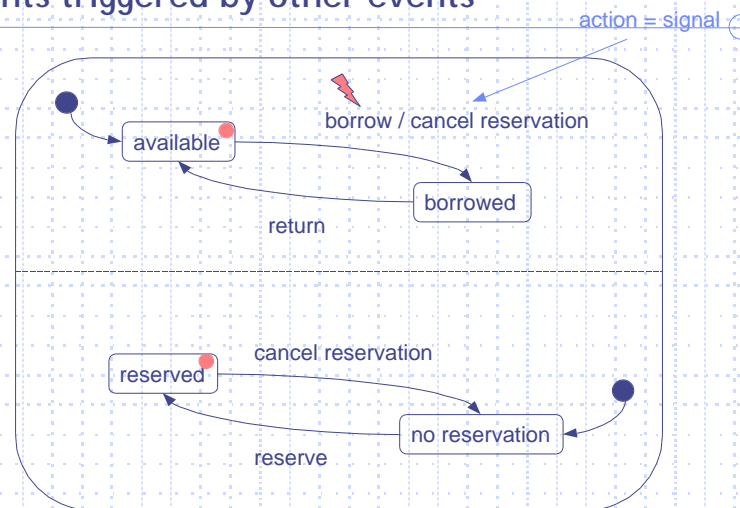
## In concurrent composite states



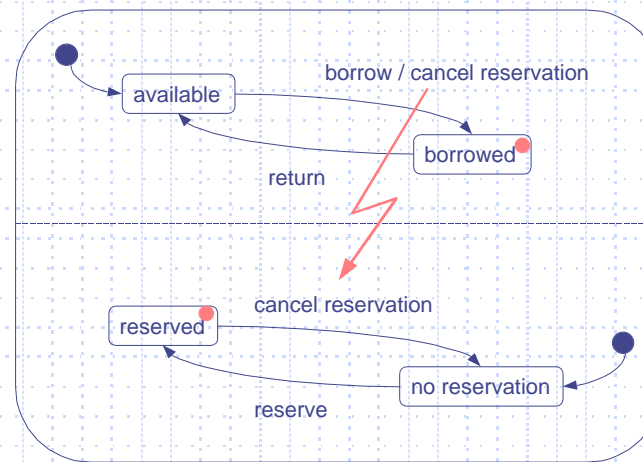
## In concurrent composite states



## Events triggered by other events



## Events triggered by other events



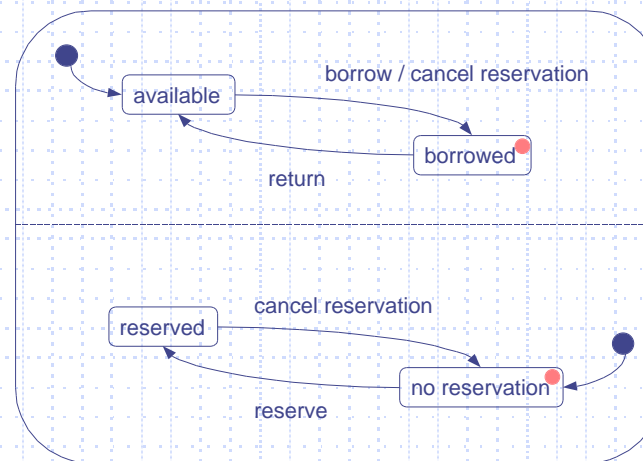
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## Events triggered by other events



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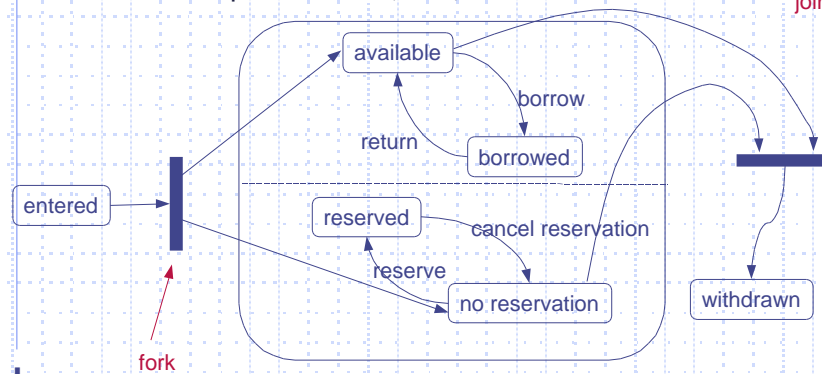
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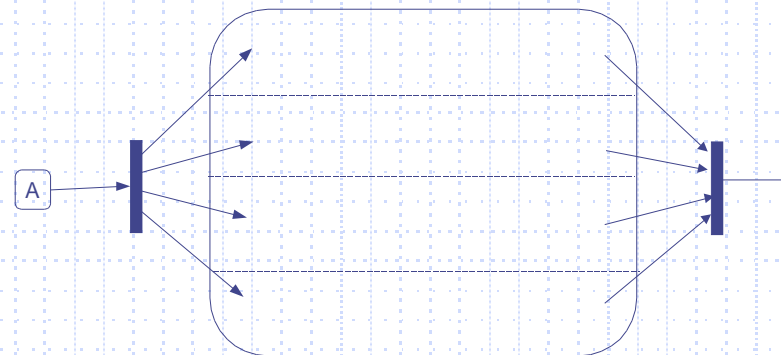
## Complex transitions

- ❖ Create parallelism
- ❖ Terminate parallelism (wait)



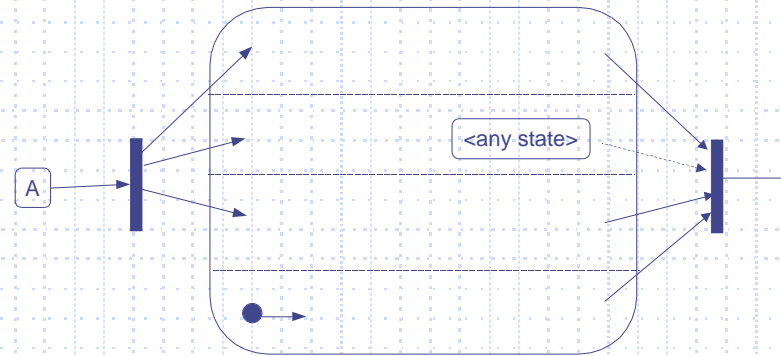
## Formation rules for complex transitions

- ❖ Enter all concurrent subregions
- ❖ Leave all concurrent subregions



## Default rules for complex transitions

- ❖ Activate the initial state.
- ❖ Leave the active substate.



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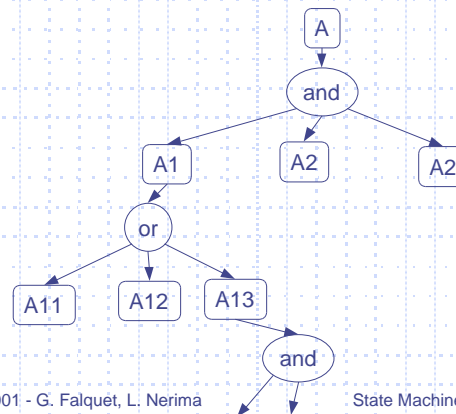
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## Structured design of state machines

- ❖ Composition of and/or states (refinement)
- ❖ => no synchronisation problems (deadlocks, etc.)
- ❖ => expressiveness slightly limited



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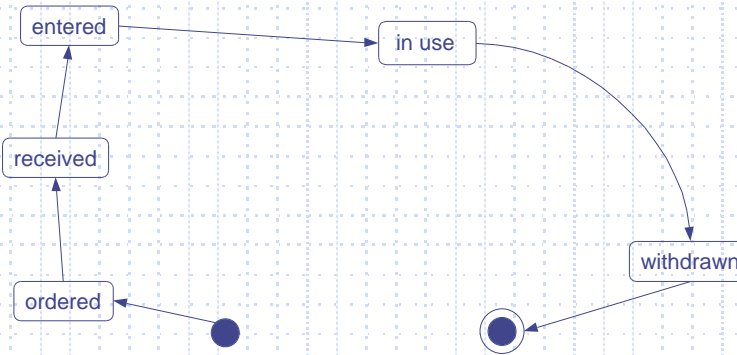
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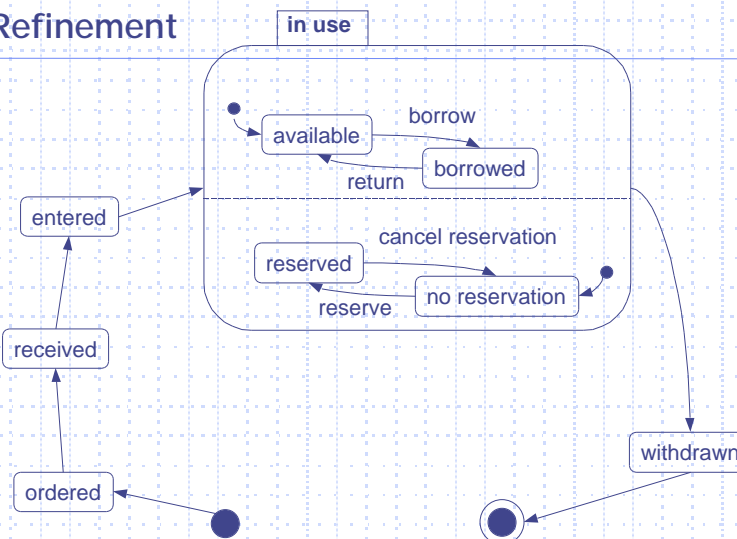
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## Example



## Refinement



## Design consideration

- ❖ State machines form the highest level of abstraction in the "dynamics" dimension
- ❖ State machines are not flow charts !
  - ❖ do not try to express algorithms, methods, computations, etc. with StM
- ❖ States must correspond to specific behaviour, conditions, etc. (avoid infinite modelling)
- ❖ Local dynamics: state machine of an object
- ❖ Global dynamics: all the state machines (with signals)

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## Implementation

- ❖ State machines are executable
  - ❖ => implementation with a state monitor/controller
  - ❖ (transaction monitors, real-time systems, ...)
- ❖ Generally: transformation into data structures and program code
- ❖ State ---> attribute value or link.
  - ❖ Book:borrowed == linked to a Loan object.
  - ❖ Report:approved == status = 'a'
- ❖ Transitions ---> execution of an operation/method.

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## State Machines and Use Cases (Douglass 2000)

- ❖ Can represent all possible scenarios on a single diagram.
  - ❖ A scenario is a path through the state machine.
- ❖ Useful to elaborate complex protocols of actor-system interaction.
  - ❖ represent actor -> system messages as triggering events and conditions
  - ❖ represent system -> actor messages as actions
- ❖ Possible to execute (simulate) the state machine to check accuracy and completeness.

<http://iamwww.unibe.ch/CHOOSE/Events/forum2k/douglass.pdf>

## Example

Interaction: The operator can enter commands to control a telescope system, subject to a number of constraints

- ❖ When the telescope is idle, the system may be configured, maintained, or commanded to move.
- ❖ The system will not accept a command when the telescope is currently moving except stop or turn off.
- ❖ Whenever the telescope is moving, monitored position is displayed in a blinking form.
- ❖ When the telescope is stopped, the monitored position is non-blinking.
- ❖ After configuration is complete, the user must reinitialise the system.
- ❖ Telescope position is displayed on a user-defined periodic basis.

